

ZS3:

Marrying Static Analyzers and Constraint Solvers to Parallelize Loops in Managed Runtimes

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“A first of its kind loop parallelizer for Java programs that combines constraint solving and static analysis to mark parallelizable loops for heterogeneous architectures using TornadoVM.”

Motivation

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```
1 public void saxpy(float alpha, float[] x, float[] y) {  
2     for(int i = 0; i < y.length; i++) {  
3         y[i] = alpha * x[i];  
4     }  
5 }
```

Simple scalar multiplication

Motivation

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```
1 public void saxpy(float alpha, float[] x, float[] y) {  
2     for(@Parallel int i = 0; i < y.length; i++) {  
3         y[i] = alpha * x[i];  
4     }  
5 }
```

Simple scalar multiplication
Parallelizable, and easy to manually annotate!

Motivation

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```
1 public void kernelThree(int nx, int ny, float[] ex, float[] hz, float[] ey) {  
2     for(int i = 0; i < nx - 1; i++) {  
3         for(int j = 0; i < ny - 1; j++) {  
4             hz[i*nx+j]=(float)(hz[i*nx+j]-0.7*(ex[i*nx+(j+1)]-ex[i*nx+j]+ey[(i+1)*nx+j]-ey[i*nx+j]));  
5         }  
6     }
```

A convolutional kernel function
What about this?

Motivation

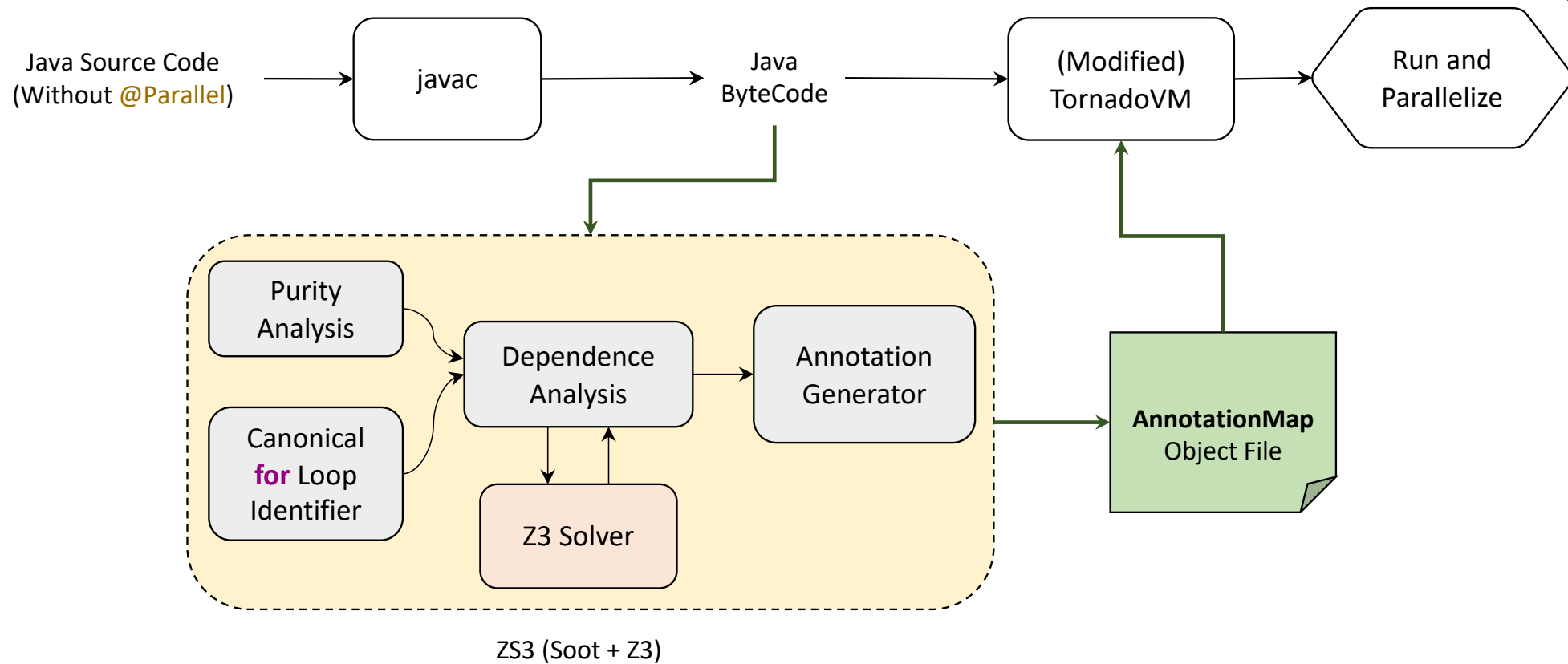
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```
1 public void foo(int[] ar) {  
2     int n = ar.length;  
3     for(int i = 0; i < n; i++) {  
4         ar[i] = bar(ar, i);  
5     }  
6 }
```

Function calls
And this?

The System

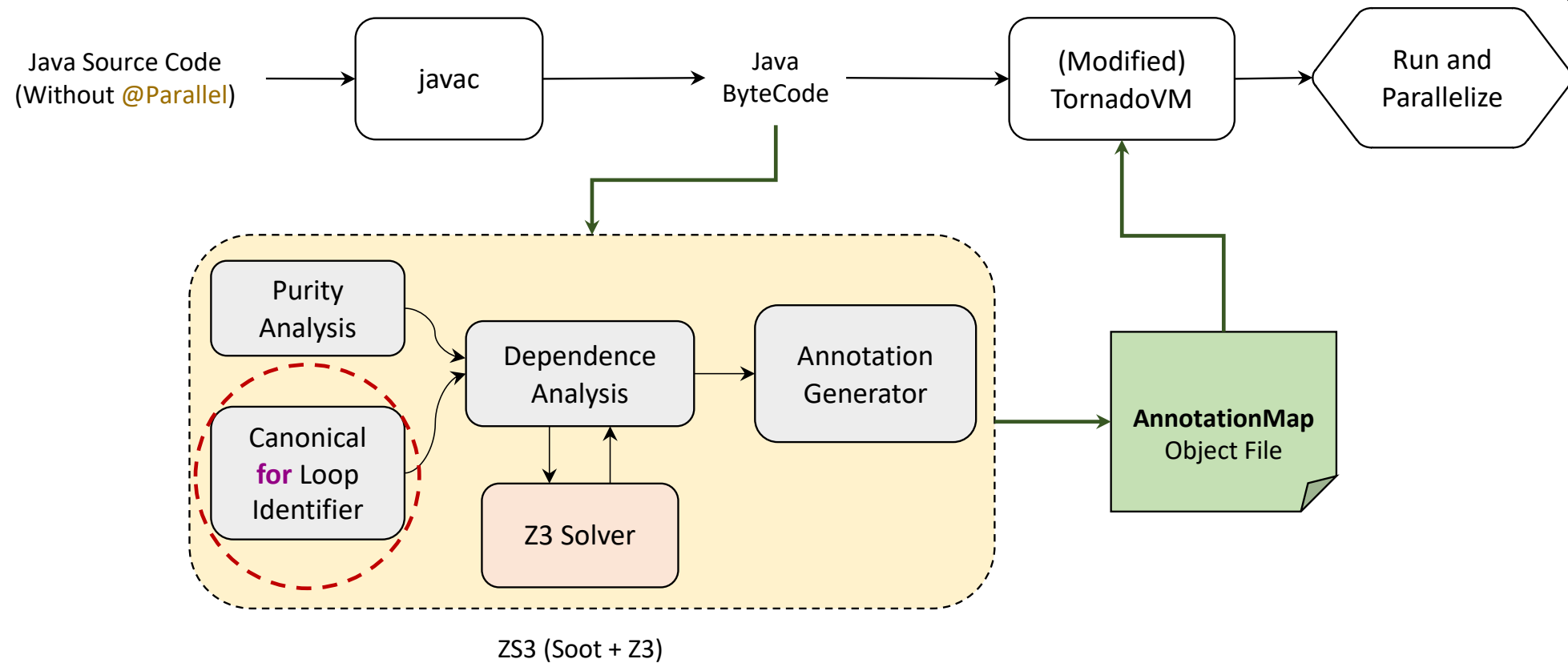
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ZS3 Architecture

The System

ZS3



The System

Canonical **for** loop identifier

```
1 public void saxpy(float alpha, float[] x, float[] y) {  
2   for(int i = 0; i < y.length; i++) {  
3     y[i] = alpha * x[i];  
4   }  
5 }
```



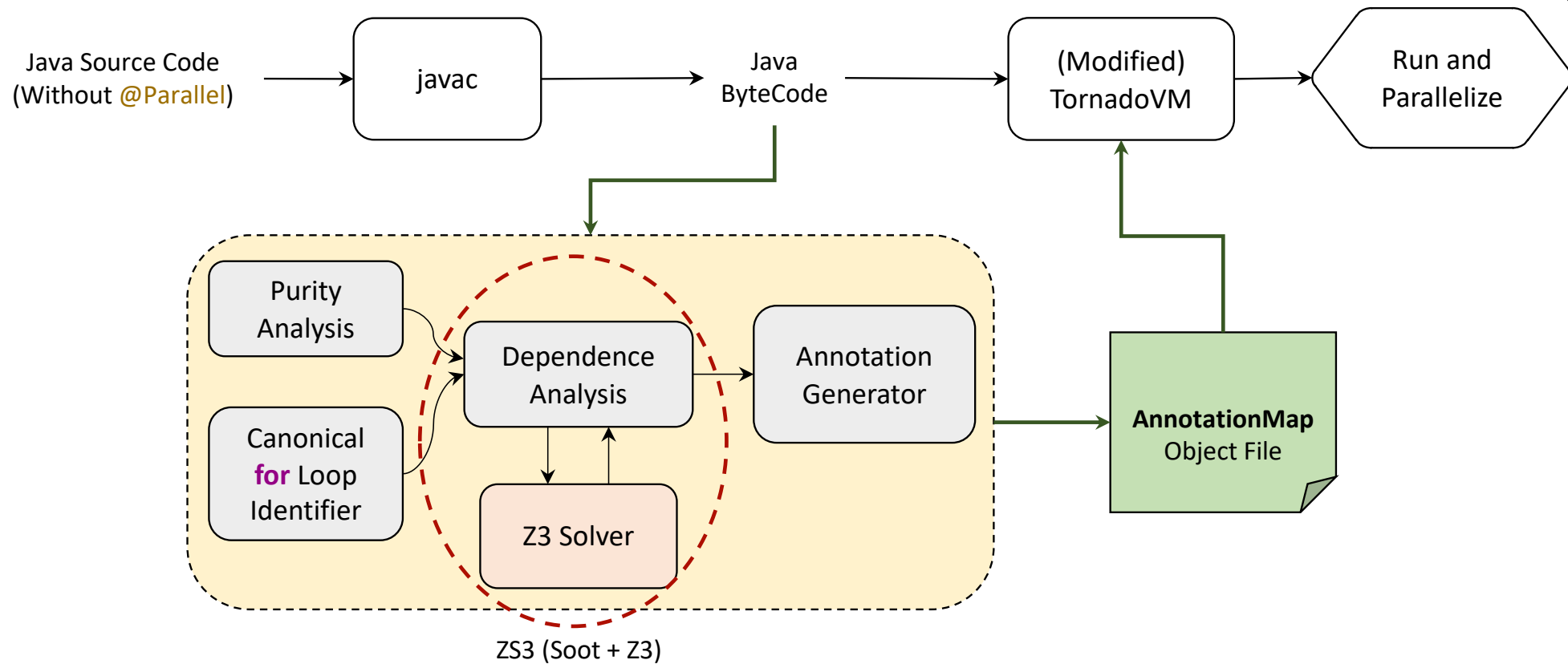
```
1 public static void saxpy(float, float[], float[])  
2 {  
3   float[] y, x; int $stack4, i;  
4   float alpha, $stack5, $stack6;  
5   alpha := @parameter0: float;  
6   x := @parameter1: float[]; y := @parameter2: float[];  
7   i = 0;  
8   label1:  
9   $stack4 = lengthof y;  
10  if i >= $stack4 goto label2;           // head  
11  $stack5 = x[i];                        // body begins  
12  $stack6 = alpha * $stack5;  
13  y[i] = $stack6;  
14  i = i + 1;  
15  goto label1;                          // body ends  
16  label2:  
17  return;                               // loop exit  
18 }
```

Java

Jimple

The System

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The System

Dependence Analysis

Scalar and Field References

- Writes to field-refs = **Dependence!**
- Writes to non-local variables = **Dependence!**
- Writes to local variables = **No dependence!**
- Reads to * = **No dependence!**

Need
variable
scopes

The System

Variable Scoping

Start	Length	Slot	Name	Signature
2	20	3	i	I
0	23	0	alpha	F
0	23	1	x	[F
0	23	2	y	[F

Compilation: `javac -g Saxpy.java`

View: `javap -p -v Saxpy`

```

1 public static void saxpy(float, float[], float[])
2 {
3     float[] y, x; int $stack4, i;
4     float alpha, $stack5, $stack6;
5     alpha := @parameter0: float;
6     x := @parameter1: float[]; y := @parameter2: float[];
7     i = 0;
8     label1:
9     $stack4 = lengthof y;
10    if i >= $stack4 goto label2;           // head
11    $stack5 = x[i];                        // body begins
12    $stack6 = alpha * $stack5;
13    y[i] = $stack6;
14    i = i + 1;
15    goto label1;                           // body ends
16    label2:
17    return;                                // loop exit
18 }
```

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Dependence Analysis

Array Indexing

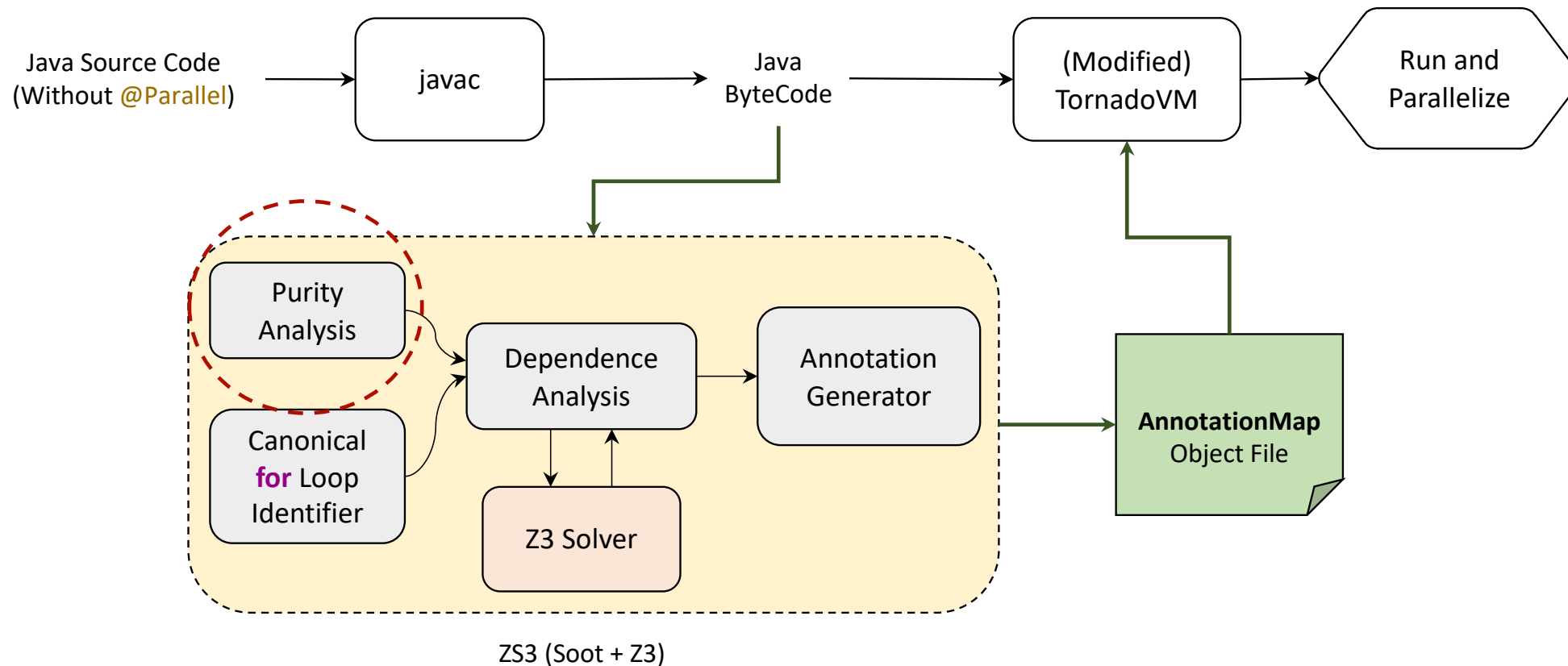
- Multiple array references in loop:
[PointsToAnalysis](#).
- For each array-write, check the dependence with all aliasing array-write and array-read.
- Using Z3! Encoding the program into logic:
[Def Chains](#).

```
1 public void foo(int ar[]) {  
2   for(int i=0; i<10000; i++) {  
3     int k1 = f1(i);  
4     int k2 = f2(i, k1);  
5     int k3 = f3(i, k2);  
6     ar[k3] = k2;  
7   }  
8 }
```


$$\begin{aligned} & (k3^u == f3(i^u, k2^u)) \wedge (k2^u == f2(i^u, k1^u)) \wedge (k1^u == f1(i^u)) \wedge \\ & \quad (i^u \geq 0) \wedge (i^u < 10000) \wedge \\ & (k3^v == f3(i^v, k2^v)) \wedge (k2^v == f2(i^v, k1^v)) \wedge (k1^v == f1(i^v)) \wedge \\ & \quad (i^v \geq 0) \wedge (i^v < 10000) \wedge \\ & \quad (i^u \neq i^v) \wedge (k3^u == k3^v) \end{aligned}$$

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The System

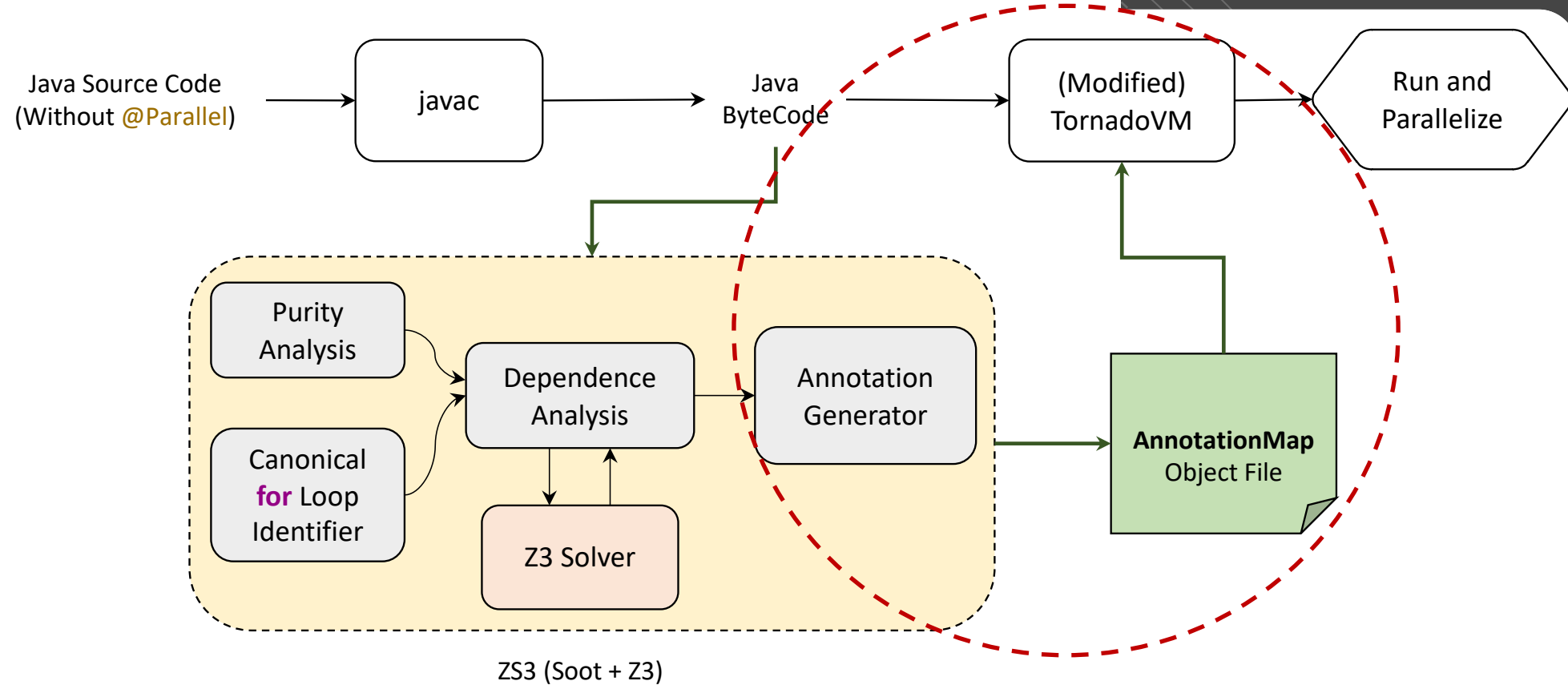
Purity Analysis

A method is pure if it does not mutate any location that exists in the program state right before method invocation

- Impure function call \Rightarrow **Non-parallelizable**.
- Sources of impurity:
 - Static field references
 - References to pre-existing objects
 - Impure function calls

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Evaluation

ZS3

Research Question 1

How many of manually parallelized loops are marked as parallelizable by ZS3?

	Marked Parallelizable	Marked Non-Parallelizable
Parallelizable	61.3%	38.7%
Non-Parallelizable	0%	100%

No False-Positives!

Evaluation

ZS3

Research Question 2

Are the overheads of static-analysis, those of storing the *AnnotationMap*, and the time spent in the VM significant?

Name	Analysis Time (s)	Class File Size (B)	AnnotationMap Size (B)	Size Overhead %	Parallel Runtime (ms)	AnnotationMap Read Time (ms)
GeoMean	1.22	4647.22	516.09	11.43	479.48	11.44
Max	4	7609.00	886.00	15.84	1546.23	25.99

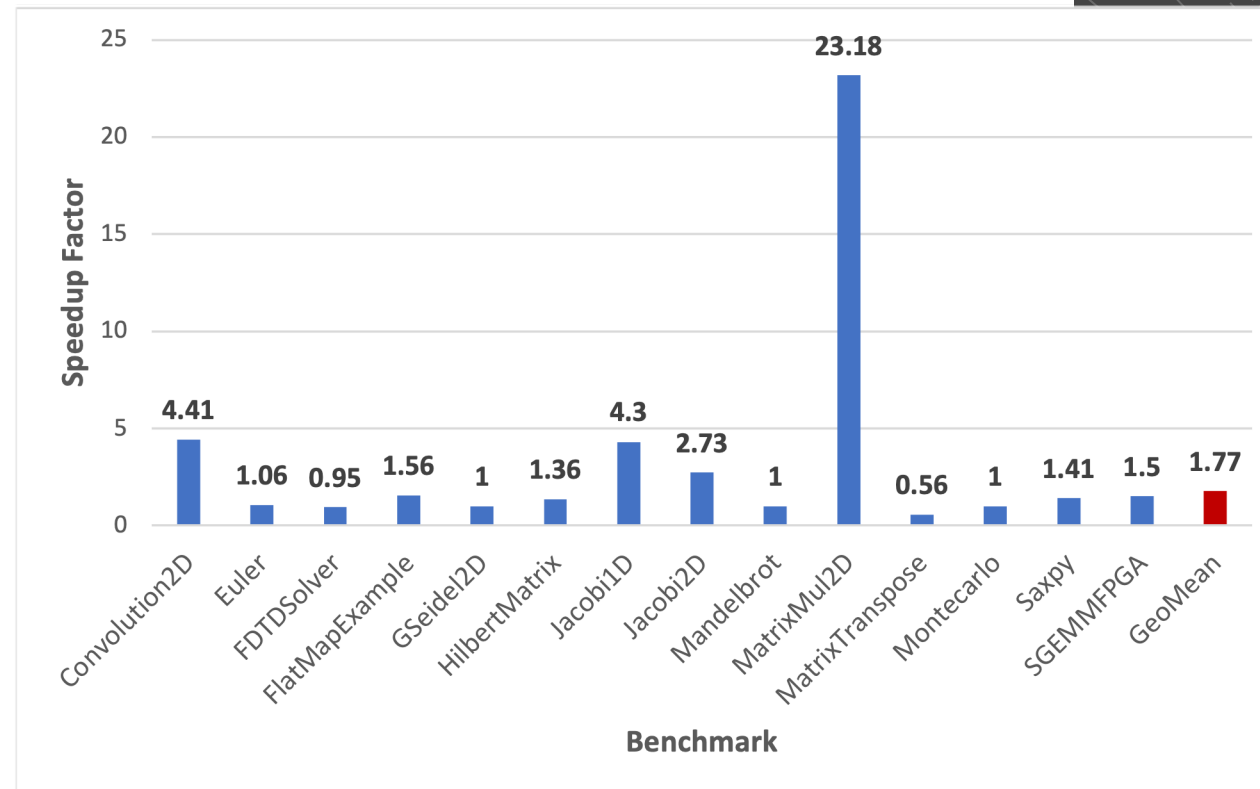
Negligible overheads

Evaluation

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Research Question 3

How good are the speedups of ZS3-marked parallel loops?



Research Question 4

What are the challenges yet to be handled by future static-analysis guided loop parallelizers?

```
1 void h(float[] output,int rows,int cols) {  
2   for (int i = 0; i < rows; i++) {  
3     for (int j = 0; j < cols; j++) {  
4       output[i*rows+j] = 1.0/((i+1)+(j+1)-1);  
5     }  
6   }  
7 }
```

when $j \geq \text{rows}$, for $i^u = 0$ and $i^v = 1$

Takeaways

ZS3

- First of its kind loop parallelizer for managed runtimes.
- Integration of Z3 with Soot \Rightarrow New possibilities.
- Precision can still be improved by integrating other techniques.

Thank You